Discovering Deterministic Finite State Automata from Event Logs for Business Process Analysis

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Highlights

- Deterministic Finite State Automata (DFAs) are employed to perform formal reasoning tasks in Process Mining [1].
- We enable the automated discovery of DFAs from event logs.
- Novel process mining quality metrics tailored to DFAs and negative examples are introduced.

Process Mining

- Process Mining [2] (PM): research area from Business Process Management.
- It analyzes process data recorded in event logs to gain insight into business processes.

DFA for PM

- Process Discovery [3],
- Conformance Checking [4],
- Compliance Monitoring [5].

Model Learning

- Active Learning, e.g. L*.
- Passive Learning, e.g.
- MDL, for positive examples,
- RPNI and EDSM, for both positive and negative examples.

- $S = \{-\} \cup S_1 \cup S_2 \cup S_3 \cup S_4$, where: • "-" is a special state;
- $S_1 = \{ \sigma \in \beta : |\sigma| \le k \}$ is the set of traces of β having length up to k;
- $\sigma[i,k]$ denoting the k-length subtrace of σ

- starting at position i, is the set of k-length prefixes of some trace in β (with length greater than k;
- $S_3 = \{\sigma[i,k] : \sigma \in \beta, |\sigma| > k, i = |\sigma| k + 1\}$ is the set of k-length suffixes of some trace in β (with length greater than k);
- $|\sigma| k + 1$ is the set of k-length subtraces of some trace in β (with length greater than k), excluding prefixes and suffixes;

by taking C

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Markovian Abstraction

- A k-th order Markovian Abstraction [6] (M^k -abstraction, for short) over a set β of traces is a finite graph $M^k = (S, E)$, with nodes S and edges $E \subseteq S \times S$, such that:
- $S_2 = \{\sigma[i, k] : \sigma \in \beta, i = 1, |\sigma| > k\}, \text{ with }$

- $E = \{(-,\sigma) : \sigma \in S_1 \cup S_2\} \cup \{(\sigma,-) : \sigma \in S_1 \cup S_2\}$
- $S_1 \cup S_3 \} \cup \{(\sigma, \sigma') : \sigma, \sigma' \in S_2 \cup S_3 \cup S_4, \exists \hat{\sigma}, i \text{ s.t. } \hat{\sigma} \in S_2 \cup S_3 \cup S_4, \exists \hat{\sigma}, i \text{ s.t. } \hat{\sigma} \in S_2 \cup S_3 \cup S_4, \exists \hat{\sigma}, i \text{ s.t. } \hat{\sigma} \in S_2 \cup S_3 \cup S_4, \exists \hat{\sigma}, i \text{ s.t. } \hat{\sigma} \in S_2 \cup S_3 \cup S_4, \exists \hat{\sigma}, i \text{ s.t. } \hat{\sigma} \in S_2 \cup S_3 \cup S_4, \exists \hat{\sigma}, i \text{ s.t. } \hat{\sigma} \in S_2 \cup S_3 \cup S_4, \exists \hat{\sigma}, i \text{ s.t. } \hat{\sigma} \in S_2 \cup S_3 \cup S_4, \exists \hat{\sigma}, i \text{ s.t. } \hat{\sigma} \in S_2 \cup S_3 \cup S_4, \exists \hat{\sigma}, i \text{ s.t. } \hat{\sigma} \in S_2 \cup S_3 \cup S_4, \exists \hat{\sigma}, i \text{ s.t. } \hat{\sigma} \in S_2 \cup S_3 \cup S_4, \exists \hat{\sigma}, i \text{ s.t. } \hat{\sigma} \in S_2 \cup S_3 \cup S_4, \exists \hat{\sigma}, i \text{ s.t. } \hat{\sigma} \in S_2 \cup S_3 \cup S_4, \exists \hat{\sigma}, i \text{ s.t. } \hat{\sigma} \in S_2 \cup S_3 \cup S_4, \exists \hat{\sigma}, i \text{ s.t. } \hat{\sigma} \in S_2 \cup S_3 \cup S_4, \exists \hat{\sigma}, i \text{ s.t. } \hat{\sigma} \in S_2 \cup S_3 \cup S_4, \exists \hat{\sigma}, i \text{ s.t. } \hat{\sigma} \in S_2 \cup S_3 \cup S_4, \exists \hat{\sigma}, i \text{ s.t. } \hat{\sigma} \in S_2 \cup S_3 \cup S_4, \exists \hat{\sigma}, i \text{ s.t. } \hat{\sigma} \in S_2 \cup S_3 \cup S_4, \exists \hat{\sigma}, i \text{ s.t. } \hat{\sigma} \in S_2 \cup S_3 \cup S_4, \exists \hat{\sigma}, i \text{ s.t. } \hat{\sigma} \in S_2 \cup S_3 \cup S_4, \exists \hat{\sigma}, i \text{ s.t. } \hat{\sigma} \in S_2 \cup S_3 \cup S_4, \exists \hat{\sigma}, i \text{ s.t. } \hat{\sigma} \in S_2 \cup S_3 \cup S_4, \exists \hat{\sigma}, i \text{ s.t. } \hat{\sigma} \in S_4, \forall \hat{\sigma} \in S_4,$ $|\hat{\sigma}| > k, 1 \le i \le |\hat{\sigma}| - k, \sigma = \hat{\sigma}[i,k], \sigma' = \hat{\sigma}[i+1,k] \}.$

Precision

Given a log ℓ and a DFA $G_{\mathcal{M}}$, let $M_{\ell}^k = (S_{\ell}, E_{\ell})$ and $M_{G_{\mathcal{M}}}^{k} = (S_{G_{\mathcal{M}}}, E_{G_{\mathcal{M}}})$ be their respective M^k-abstractions, C the Levenshtein-distance-based cost matrix, and let μ_C : $E_{G_{\mathcal{M}}} \rightarrow E_{\ell}$ be a partial function, solution of the assignment problem represented by C. The (Markovian-abstractionbased) k-th order precision of $G_{\mathcal{M}}$ wrt ℓ is defined as:

$$AAP^{k}(\ell, G_{\mathcal{M}}) = 1 - \frac{\sum_{e \in E_{G_{\mathcal{M}}}} C(e, \mu_{C}(e))}{|E_{G_{\mathcal{M}}}|},$$

$$C(e, \mu_{C}(e)) = 1, \text{ if } \mu_{C}(e) \text{ is undefined.}$$

Given a log ℓ and a DFA $G_{\mathcal{M}}$, let $M_{\ell}^k = (S_{\ell}, E_{\ell})$ and $M_{G_{\mathcal{M}}}^{k} = (S_{G_{\mathcal{M}}}, E_{G_{\mathcal{M}}})$ be their respective M^k-abstractions, C the Boolean cost matrix, and let $\mu_C: E_\ell \to E_{G_M}$ be a partial function, solution of the assignment problem represented by C. The (Markovian-abstraction-based) k-th order fitness of $G_{\mathcal{M}}$ wrt ℓ is defined as:

$$MAF^k(\ell, G_{\mathcal{M}}) = 1$$

where F_e stands for the frequency of edge e in E_{ℓ} and taking $C(e, \mu_C(e)) = 1$, if $\mu_C(e)$ is undefined.

Conclusions

- Active learning algorithms are not suitable to generate DFAs from real-life event logs.
- Declare Miner and passive learning algorithms construct
- Passive learning algorithms generate simpler DFAs than Declare Miner.

Future Work

- Learn LTL_f formulae:
- directly from logs, or,
- going through Alternating Finite Automa.
- both approaches possible with SAT or ASP techniques.

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Fitness

 $\Sigma_{e \in E_{\ell}} C(e, \mu_C(e)) F_e$ $\sum_{e \in E_{\ell}} F_e$

DFAs with similar values of generalization and precision.

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